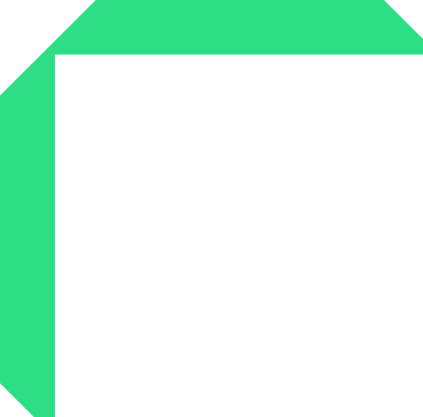




100 BC

Meeting C++ 2025-11-06

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About Marc Mutz

- Principal Software Engineer at The Qt Company
- For the purposes of this talk:
 - The guy that annoys other Qt devs in API / ABI reviews
 - In particular, not an expert in this topic
 - just more experienced than most devs in these matters

About This Talk

- If you need definite answers, talk to your vendor
- Out of Scope are:
 - C++20 Modules
 - Linker Scripts
 - Static Linking
 - Combining different compiler-vendors, -versions, or -flags

Overview

- Motivating Example
- Definitions
- A look at the KDE BC Page
- Techniques
- Cheating
- Q & A

Example

- Lib v1.0

```
int scan(int fd);
```

- Lib v1.1:

```
enum class Options { ~~~ };
```

```
int scan(int fd, Options opt = {});
```

Example

- Lib v1.1:

```
enum class Options { ~~~ };
```

```
int scan(int fd, Options opt = {}); // BiC
```

- Lib v1.1p1:

```
int scan(int fd);
```

```
int scan(int fd, Options opt);
```

Example

- Lib v1.1p2:

```
#if LIB_REMOVED_SINCE(1, 1)
```

```
int scan(int fd);
```

```
#endif
```

```
int scan(int fd, Options opt = {});
```

Definitions

Definitions

- Qt / KDE
- BC / SC
- ABI / API
- forwards / backwards xC
- name mangling / exported symbols

Definitions – Qt and KDE

- Qt: A cross-platform application development framework written in C++
- KDE: A (Unix) Desktop Environment based on Qt
- Both promise BC and SC within major release cycles

Definitions — BC

“A library is binary compatible, if a program linked dynamically to a former version of the library continues running with newer versions of the library without the need to recompile.”

— KDE BC wiki

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Observations — BC

- BC is not defined (by KDE) for static linking
 - Doesn't stop people from trying, though
- "Continues running" is vague
 - "Still links" is just one (necessary) condition
 - Not sufficient. Also need "behaviour compat".
 - Even a bugfix may break programs



Observations — BC

- You don't need a library to talk about BC
 - A type can be BC
 - A function can be BC
- A library is just the unit of shipment
 - BC is a property of smaller units

Definition — SC

"If a program needs to be recompiled to run with a new version of a library but doesn't require any further modifications, the library is *source compatible*."

— KDE BC wiki

Observations — SC

- Inline library code can be SiC, but usually not BiC
 - Compiled into the library user
 - Doesn't change when the library changes
- BC != SC
 - Adding a defaulted argument to a function: SC && !BC
 - Adding a new overload of a function: BC && !SC

Definitions — ABI

- Application Binary Interface
 - “what the linker sees”
 - “below C++”
 - “mangled names”
 - “exported symbols”

Definitions — API

- Application Programming Interface
 - “what the compiler sees”
 - “defined by C++”
 - “overloading” / “default arguments”
 - “typedef”

Definitions — backwards xC

- “Normal” compat:
 - Code compiled against *old* library running against *new* library
- Qt provides backwards BC / SC¹ in major releases
 - 5.0 → 5.15
 - 6.0 → 6.???

¹ see [QUIP-6](#) for acceptable SC breaks

Definitions — forwards xC

- “Reversed” compat:
 - Code compiled against *new* library running against *old* library
- Qt provides forwards BC / SC in minor releases
 - 5.15.0 ← 5.15.19
 - 6.8.0 ← 6.8.???

Definitions — Name Mangling

- Creates unique names for overloaded functions
- Encodes things you can overload on:
 - Name
 - incl. any scopes ("FQN")
 - signature (= number and type of arguments)
 - Incl. template arguments
 - Incl. cv- and l/rvalue qualifiers of member functions

Definitions — Name Mangling

- Does *not* encode:
 - Argument names and defaults, if any
 - Argument top-level cv-qualifiers (except SunCC in its days)
 - noexcept
 - Access specifier (except MSVC)
 - Return type (except MSVC)

Definitions — Name Mangling

Conclusion:

- Change to name mangling inputs = ABI break
- Change to non-inputs = BC
- Platform differences!

Definitions — Exported Symbols: C++

- extern linkage
 - other TUs can access
 - needs declaration
 - duplicates not allowed
- static linkage
 - only this TU can access
 - duplicates remain
- inline linkage
 - other TUs can access
 - needs definition
 - duplicates are merged

Definitions — Exported Symbols: DSO

- Unix: callable from outside .so/.dylib by default
 - inline (code copied into caller)
 - extern (library code called)
- Windows: callable from outside .dll by default
 - inline (code copied into caller)

Definitions — Exported Symbols

- Exporting is reducing (Unix) / extending (Win) the set of symbols callable from outside the DSO
 - Windows: `__declspec(dllexport)`
 - Unix: `-fvisibility=hidden +
__attribute__((visibility("default")))`

Definitions — Exported Symbols

- These attributes need to differ between building the library and building against the library
- Hidden behind macros
- Qt: `QT_(MODULE)_EXPORT`
 - generated by CMake

Definitions — Exported Symbols

- Types don't need exporting (are not symbols)
 - “Exported Class” is just “export all members”
 - `class QT_CORE_EXPORT QObject { ~~~ };`
- Variables and function do
 - incl. vtable of polymorphic classes (unless inline)

Definitions — Exported Symbols

- “Wholesale” exporting “the class” exports:
 - all member functions
 - all static data members (incl. vtable)
 - inline functions also get exported(!)
 - and, on Windows, the DLL-exported version called
 - or not, if the optimizer inlines it `\_(ツ)_/`

Definitions — Exported Symbols

- “Wholesale” exporting “the class” is recursive:
 - nested classes are “exported”, too

Definitions — Exported Symbols

- “Wholesale” exporting “the class” does not export:
 - (hidden) friend functions
 - non-static data members (not variables)
 - member function templates
 - nested class templates

Definitions — Exported Symbols

Conclusion:

- You don't want to export inline symbols
- You'll want to *selectively* export extern symbols
 - those that need to be called from outside the DSO
- You don't want to export an "inline" vtable



"Golden Rules" Of Class Exporting

- Export polymorphic classes "wholesale"
 - and define the dtor out-of-line (even if empty / = default)
- Do not "wholesale" export non-polymorphic classes
 - Export individual functions / variables only
 - e.g. not inlines, not privates (unless called from inline code)

Definitions — Exported Symbols

- Exported symbols are part of the ABI
 - incl. exported inlines (!)
- Non-exported symbols are not
 - the fewer exported symbols, the fewer BC headaches
 - exporting allows “library-private” symbols

Note On Unexported Inline Symbols

- Linker deduplicates them only *per executable*
 - each DLL, and the app, each get their own copy
 - harmless for functions
 - except their addresses differ
 - surprising for variables

Note On Unexported Inline Symbols

- Solution: Aggressive DRY / SCARY
 - iow: extract such things into separate functions
 - “pull” these functions “behind the ABI boundary”
 - de-inline and export

A Look at the KDE BC Page

- There are many do's and don't's that we have no time for today, let's peek at the page together:



Techniques

Techniques

- plmpl
- reserved fields
- move semantics
- virtual functions / hooks
- inline namespaces

Techniques – pImpl

- “Pointer to implementation” / “Compiler Firewall”
 - `class Private; Private* pImpl; // only data member`
 - externalizes all private state (and functions)
 - private parts can vary independent of public iface

Techniques – pImpl

- Field new'ed up in ctors
 - memory allocation (or pooling)
 - mitigation: implicit sharing, base class reuse, ...
 - Almost all member functions must be out-of-line
 - Exceptions:
 - move SMFs and swap()

Techniques – plmpl

- Lots of gotchas:
 - shallow const; back pointers; Base::plmpl reuse in Derived
- See article series on [Heise](#) (DE) / [-Wmarc](#) (EN)



Techniques – plmpl

Conclusion:

- Effectiveness: ++
- Runtime overhead: --
 - Moves and swaps are cheap, though

Techniques – reserved fields

- Directly-embedded unused data members
 - `[[maybe_unused]] void* reserved = nullptr;`
 - unused bits in bit-fields
 - `[[maybe_unused]] uint reserved = 0U;`
 - explicit padding: `[[...]] char padding_n[M] = {};`

Techniques – reserved fields

- Fields need to be init'ed, even if unused
 - use NSDMI: = 0U/nullptr
 - if inline code inits, field is already severely restricted
 - e.g. 0 needs to be a valid value (or partially-formed)
 - if inline code destroys, field must be trivially destructible

Techniques – reserved fields

- Construction and destruction must be out-of-line!
 - copy SMFs, dtor, any custom ctor (incl. default ctor)
- Exceptions:
 - move SMFs, swap() → inline
 - delegating constructors → may be inline

Techniques – reserved fields

Conclusion:

- Effectiveness: 0
 - Limited capacity (literally) for change
- Runtime overhead: +
 - Pay only for what you use (almost)

BC and Move Semantics

- Move SMFs should be fast → inline
- But they shouldn't break BC → out-of-line
- Conflict of interest?

BC and Move Assignment

- Natural implementation of move-assignment can be inline:
 - `swap(this->m_data, other.m_data)`
 - `plmpl` and `reserved` field alike
 - “valid, but unspecified”
- BUT: assignment doesn't destroy LHS
 - OK if type holds only memory resources
 - Alternative: `move-and-swap` (needs move ctor)

BC and Move Constructor

- Natural implementation of the move constructor can also be inline:
 - `m_data{std::exchange(other.m_data, 0U/nullptr)}`
 - “Partially-Formed” State (→ [Meeting C++ 2020](#))



BC and Move Constructor

- BUT: doesn't work with smart_ptr
 - Compiler wants to call data member dtor
 - Doesn't need it, but breaks encapsulation!
 - @ ISO C++: can we fix this?
 - noexcept move ctors oughtn't call data member dtors

BC and Move Constructor

- smart_ptr plmpl work-around
 - Make the smart_ptr dtor call an out-of-line function
 - Idea: call isn't emitted from move ctor
 - Inside the library, compiler can inline said function
 - No performance overhead
 - BUT: MSVC LTCG doesn't play ball (see references)

BC and Move Semantics

Conclusion:

- Move SMFs and `swap()` can (and ought to) be inline
- ISO C++ and / or MSVC may have some bugfixing to do

Techniques – Virtual Hook

- Problem: can't add virtual function to a class
- Qt/QObject-specific work-around:
 - non-virtual function calls impl slot through meta-object
 - works, because the call is delivered via virtual function
 - idea can be generalized, though

Techniques – Virtual Hook

- Needs to have been thought of in advance:

protected:

```
virtual void virtual_hook(int id, void *args) const;
```

- Must have be overridden in every class

Techniques – Virtual Hook

- New function:

```
int pseudoVirtual(int arg) {  
    int argv[2] = {0, arg};  
    virtual_hook(ID, argv);  
    return argv[0];  
}
```

- Impl in virtual_hook():

```
if (id == ID) {  
    auto argv =  
        static_cast<int*>(args);  
    argv[0] = op(argv[1]);  
}
```

Techniques — Virtual Hook

Conclusion:

- Good for isolated use
- Only for types that are already polymorphic
- Needs to be added ante hoc
- Hard-to-automate task when extending hierarchy

Techniques — Inline Namespaces

- Type versioning at the C++ level
- Lib v1.0:
 - `struct S { int x, y, z; }; int foo(const S &);`
- Lib v1.1:
 - `struct S { int x, y, z; Flags f; }; int foo(const S &);`

Techniques – Inline Namespaces

```
inline namespace S_V1_0 {  
    struct S { int x, y, z; };  
}  
  
int foo(const S &s);
```

Techniques – Inline Namespaces

```
inline namespace S_V1_0 {  
    struct S { int x, y, z; };  
}  
  
int foo(const S_V1_0::S &s); // unchanged!  
  
inline namespace S_V1_1 {  
    struct S { int x, y, z; Flags f = {} };  
}  
  
int foo(const S &s); // new function!
```

Techniques — Inline Namespaces

- Empowers users:
 - “live at HEAD”: use unqualified `S`
 - “pin old API”: use `S_V1_0::S` explicitly
 - can even be mixed in the same TU

Techniques — Inline Namespaces

- Combinatorial explosion in N-ary functions, $N > 1$:
 - `replace(String &, const String &, const String &); // 8×`
- Different components extend at different times
 - need different namespaces
 - $V1_0 \rightarrow S_V1_0$
 - or lots of boilerplate
 - importing all unchanged `V1_0` types into `V1_1`

Techniques — Inline Namespaces

Conclusion:

- Good for isolated use (e.g. Config struct)
- Quickly gets out of hand when used widely
- Needs to be added ante hoc
- Lots of hard-to-automate tasks on version bumps

Cheating

Cheating

- Separating API and ABI
- Removing functions from one, but not the other
- Qt's user-config'able, rolling BC window
- Bonus: inlining out-of-line functions

Cheating — Observations

- Types don't exist in the ABI
 - only as a collection of functions acting on them
 - there is no “Point” in the ABI
 - there's only the agreement of certain functions as to type and location of Point's fields

Cheating — Observations

- ABI is maintained as long as
 - (existing) memory layout doesn't change
 - (existing) functions are not removed

Cheating — Example API vs. ABI

```
struct LIB_EXPORT Point {  
    int m_x, m_y;  
    void setX(int x);  
    int x() const;  
  
    void setY(int y);  
  
    int y() const;  
  
};
```

- `offsetof(Point::m_x) == 0`
- `offsetof(Point::m_y) == 4`
- `Point_setX_i();`
- `Point_x_v();`
- `Point_setY_i();`
- `Point_y_v();`

Cheating — Example Func Addition

```
struct LIB_EXPORT Point {  
    int m_x, m_y;  
    void setX(int x);  
    int x() const;  
  
    void setY(int y);  
  
    int y() const;  
    void setY(long long y);  
};
```

- `offsetof(Point::m_x) == 0`
- `offsetof(Point::m_y) == 4`
- `Point_setX_i();`
- `Point_x_v();`
- `Point_setY_i();`
- `Point_y_v();`
- `Point_setY_ll();`

Cheating — Example Func Addition

```
struct LIB_EXPORT Point {  
    int m_x, m_y;  
    void setX(int x);  
    int x() const;  
#if 0  
    void setY(int y);  
#endif  
    int y() const;  
    void setY(long long y);  
};
```

- `offsetof(Point::m_x) == 0`
 - `offsetof(Point::m_y) == 4`
 - `Point_setX_i();`
 - `Point_x_v();`
 - `Point_y_v();`
 - `Point_setY_ll();`
- (oops)

Cheating — Example Func Addition

```
struct LIB_EXPORT Point {  
    int m_x, m_y;  
    void setX(int x);  
    int x() const;  
#if MAGIC_REMOVED_API_ONLY  
    void setY(int y);  
#endif  
    int y() const;  
    void setY(long long y);  
};
```

- `offsetof(Point::m_x) == 0`
- `offsetof(Point::m_y) == 4`
- `Point_setX_i();`
- `Point_x_v();`
- `Point_setY_i();`
- `Point_y_v();`
- `Point_setY_ll();`

Cheating — removed_api.cpp

// point.cpp

```
void Point::setY(long long i) {  
    ~~~~  
}  
void Point::setY(int i) {  
    ~~~~  
}
```

// removed_api.cpp

```
#define MAGIC_REMOVED_API_ONLY 1  
#include "point.h"
```

Cheating — removed_api.cpp

// point.cpp

```
void Point::setY(long long i) {
```

```
    ~~~~
```

```
}
```

```
void Point::setY(int i) {
```

```
    _____
```

```
}
```

// removed_api.cpp

```
#define MAGIC_REMOVED_API_ONLY 1
```

```
#include "point.h"
```

```
void Point::setY(int i) {
```

```
    setY(static_cast<long long>(i)); // new overload
```

```
}
```

Cheating — REMOVED_SINCE

- That's all!
- You can get fancy and version the thing:
 - `API_REMOVED_SINCE(1, 1)`
- You need this macro per-library
 - so let your build system generate these
 - but only one `removed_api.cpp` per library

Cheating — REMOVED_SINCE

- more sophisticated symbol versioning exists
- REMOVED_SINCE's advantage is
 - purely C++-based
 - every C++ developer can understand how to apply it
 - cross-platform
 - works post-hoc

Cheating — Homework

- Use `REMOVED_SINCE()` to remove `noexcept` from an exported function again
 - Think about syntax in header and `removed_api.cpp`
 - Also think about semantics of removing `noexcept`

REMOVED_SINCE — Cut-Off Version

- What happens if I introduce a cut-off version X.Y?
 - i.e. `REMOVED_SINCE()` = 0 for all versions before X.Y
- I have just introduced a (user-config'able) cut-off for the ABI zombie functions in my DSO
 - IOW: a config'able, rolling BC window

A Rolling BC Window

- Upper bound:
 - The library version I'm building against
- Lower bound:
 - Said cut-off version
- Qt's called -disable-deprecated-up-to

Cheating — REMOVED_SINCE

- Biggest drawback is: doesn't work for virtuals
 - necessary precondition not met:
 - need to be able to add something to remove something else
 - if it weren't for virtuals, Qt could stay two decades more on 6.x
 - alternatively, Qt 7 could be BC with Qt 6

Key Take-Aways

1. ABI $\neq$ API; both can vary independently
2. Qt has a user-config'able, rolling BC window
3. "Wholesale" exported classes and virtual functions are big impediments to ABI evolution

Q & A

References

- KDE BC page
- Marc Mutz - On removing functions from the API, but not the ABI – Qt development ML 2022-07-13
- Marc Mutz - Partially-formed Objects for fun and profit - Meeting C++ 2020
- Marc Mutz - Pimp My Pimpl – Heise Developer 2010
 - English Translation thereof
- Discussion about MSVC LTCG bug in Qt bug tracker
 - Stack Overflow discussion: “Where does the destructor hide in this code?”
- QUIP-6: Acceptable Source-Incompatible Changes
- Marc Mutz – Fun With Exceptions - -Wmarc 2010-08-04
- Qt documentation: QT_DISABLE_DEPRECATED_UP_TO

Cheating — Bonus Slides

Cheating — Inlining Functions

- Next Level: inlining formerly out-of-line functions
- Idea:
 - REMOVE_SINCE the out-of-line function
 - add the inline function
- Two problems...

Cheating — Inlining Functions

- Obvious problem: the two don't overload
 - just use `#else`
- Not-so-obvious problem: ODR violation
 - `REMOVED_SINCE` is, too, but non-actionable
 - `inline` vs. `extern` linkage, or even vs. `exported` *is* actionable

Cheating — Inlining Functions

- We have three domains now:
 - removed_api.cpp
 - callers external to the library
 - callers internal to the library

Cheating — Inlining Functions

- external users should see the function as inline
- `removed_api.cpp` must see it as out-of-line
 - to maintain ABI
- internal users must see it as out-of-line, too
 - else: ODR violation when linked with `removed_api.cpp`

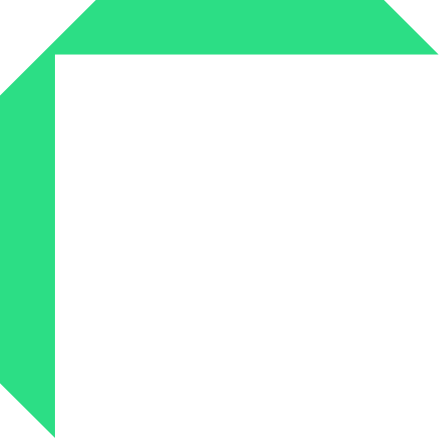
Cheating — INLINE_SINCE

```
// point.h
struct Point { ~~~
    API_INLINE_SINCE(1, 1)
    int x() const;
~~~ };
#if API_INLINE_IMPL_SINCE(1, 1)
int Point::x() const { ~~~ }
#endif
```

```
// removed_api.cpp
~~~
#include "point.h" // inlined API
```

Cheating — `INLINE_SINCE`

- `INLINE_SINCE`
 - inline
 - external callers
 - `/* not inline */`
 - `removed_api.cpp`
 - internal callers
- `INLINE_IMPL_SINCE`
 - 1
 - external callers
 - `removed_api.cpp`
 - 0
 - internal callers



The End